Logic and Constraint Programming



5- Rule-based systems

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Rule-based systems



- Rules are the main way to express knowledge in many fields of I.A.
- · Most common rules are:
 - production rules (eg.: Drools)
 - logic programs (eg.: Prolog)
- They are similar, but realized in a dual way



· Modus Ponens:

$$\frac{\langle p(x), p(X) \rightarrow q(Y) \rangle}{q(y)}$$

if it holds that p(X) implies q(Y) and p(x) holds, then q(y) holds

Es.: If it rains, then the street is wet.

Here it rains.

Then, here the street is wet.



Modus Ponens:

$$\frac{\langle p(x), p(X) \rightarrow q(Y) \rangle}{q(y)}$$

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Es.: If it rains, then the street is wet.

Here it rains.

Then, here the street is wet.

implication premise conclusion



Production rules

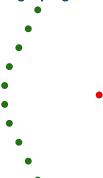
- Forward-chaining
- The facts activate rules that generate new facts
- Pattern matching
- Parallelism

- Backward-chaining
- From goal to facts, applying rules in a backward way
- Unification
- Backtracking



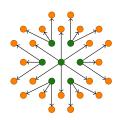
Production rules

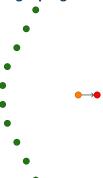






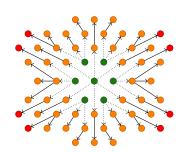
Production rules

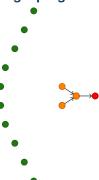






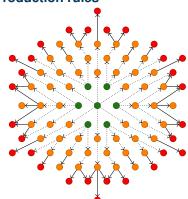
Production rules

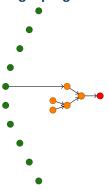




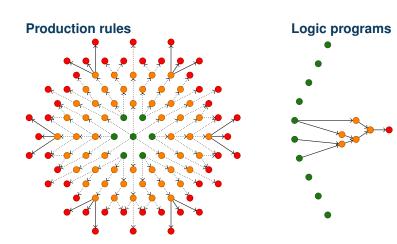


Production rules





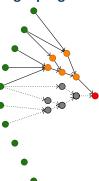






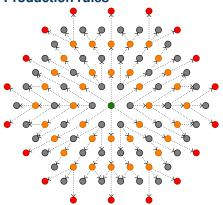
Production rules

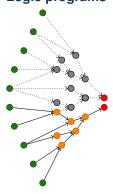






Production rules







- Production Rule Systems (PRS):
 - are Rule Based Systems (RBS),
 - · are based on the Modus Ponens principle,
 - · rely on a reactive/generative approach

WHEN A PRS IS A RIGHT CHOICE?



- The problem is too complex for traditional coding approaches: rules provide a more abstract view, preventing fragile implementations
- The problem is not fully known
- Flexibility, when system logic changes often over time
- Domain knowledge readily available

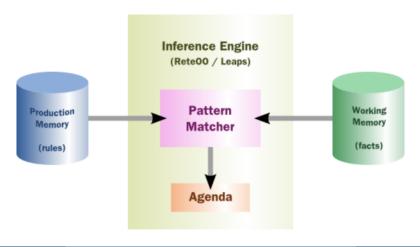
EXAMPLE SCENARIO





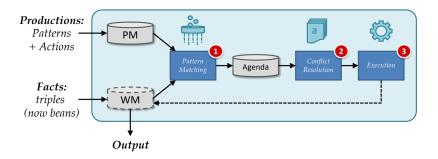


Architecture and working schema





Architecture and working schema



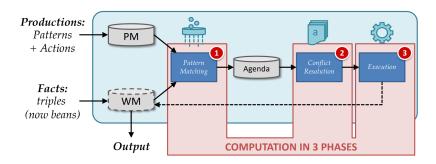


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- Rules are stored in the Production Memory (PM)
- Facts are stored in the Working Memory (WM), where they can be changed or retracted
- Inference engine applies to data in the WM the rules in in the PM to deduce new information
- The Agenda deals with the execution order in case of conflicts, using conflict resolution strategies

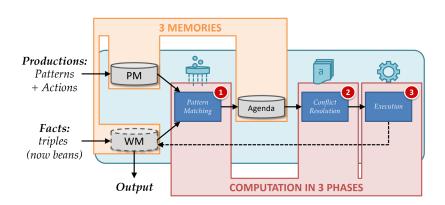


Architecture and working schema





Architecture and working schema



Pattern Matching: RETE

algorithm

RETE ALGORITHM



RETE is a **pattern matching** algorithm for implementing **rule-based systems**.

The Rete algorithm was designed by *Charles L. Forgy* of Carnegie Mellon University, first published in a working paper in 1974, and later elaborated in his 1979 Ph.D. thesis and a 1982 paper.



RETE NETWORK



The **Rete network** is the *brain* behinf the Rete algorithm

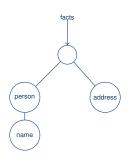
It is made of **nodes** that each hold a list of objects that **satisfy some** associated condition

The original Rete algorithm worked out of **facts**, while commercial engines have evolved to be object-oriented nowadays

RETE NETWORK >ALFA NODES



The discrimination tree starts with Alfa nodes



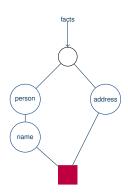
Alfa nodes are created for each fact, then **attributes** are appended

Each node represents an additional test to the series of **conditions** applied upstream

RETE NETWORK >BETA NODES



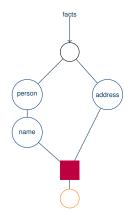
Nodes are then connected across facts into **Beta nodes**



Those nodes combine the list of facts that verify conditions on one branch with the list of facts that verify the conditions on another branch

RETE NETWORK



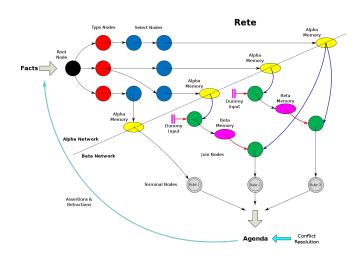


The path eventually ends with the **action part** of the rule

The content of the actions is irrelevant for the Rete network



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- The Rete implementation used in Drools is called ReteOO
- It is an enhanced and optimized implementation of the Rete algorithm specifically for object-oriented systems

RETEOO SOBJECT TYPE NODE



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When using ReteOO, the root node is where all objects (facts) enter the network. From there, it immediately goes to the **ObjectTypeNode**.

The ObjectTypeNode helps to reduce the workload of the rules engine. To make things efficient, the ObjectTypeNode is used so that the engine only passes objects to the nodes that **match the object's type**

An inserted object retrieves a list of valid ObjectTypesNodes through a lookup in a HashMap from the object's class. If this list does not exist, it scans all the ObjectTypeNodes to find valid matches



AlfaNodes are used to evaluate literal conditions. When a rule has multiple literal conditions for a single object type, they are linked together. E.g., if an application asserts an object, it must first satisfy the first literal condition before it can proceed to the next AlfaNode

AlfaNodes are propagated using ObjectTypeNodes. Each time an AlfaNode is added to an ObjectTypeNode, it adds the literal value as a key to the **HashMap** with the AlfaNode as the value.

RETEOO >ALFANODES



When a new instance enters the ObjectType node, rather than propagating to each AlfaNode, it retrieves the correct AlfaNode from the HashMap. This avoids unnecessary literal checks.

When facts enter from one side, you may do a **hash lookup** returning potentially valid candidates (referred to as indexing). At any point a valid join is found, the Tuple joins with the Object (referred to as a partial match) and then propagates to the next node.



BetaNodes are used to compare two objects and their fields. The objects may be of the same or different types

Alfa memory refers to the left input on a BetaNode. Beta memory is the term used to refer to the right input of a BetaNode

When facts enter from one side, if a valid join is found, the object (referred to as a partial match) and then propagates to the next node

RETEOO STERMINAL NODES



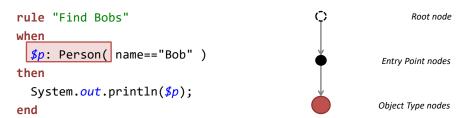
Terminal nodes are used to indicate when a single rule matches all its conditions (that is, the rule has a full match). A rule with an OR conditional disjunctive connective results in a sub-rule generation for each possible logical branch. Because of this, one rule can have **multiple terminal nodes**

Rete00 Examples

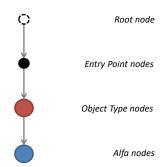
```
rule "Find Bobs"
when
    $p: Person( name=="Bob" )
then
    System.out.println($p);
end
```

Root node

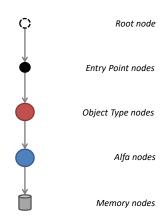
Entry Point nodes

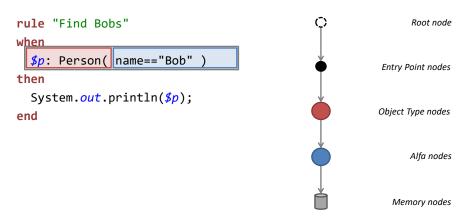


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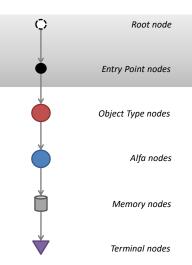


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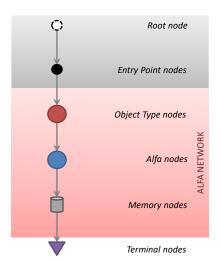




NB: facts in a (Alfa) Memory Node match with a simple pattern!

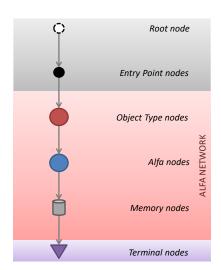


```
rule "Find Bobs"
when
$p: Person( name=="Bob" )
then
   System.out.println($p);
end
```

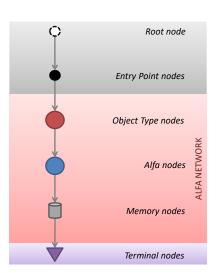


```
rule "Find Bobs"
when

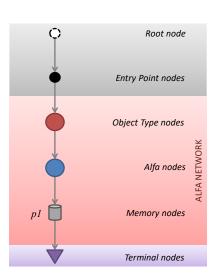
$p: Person( name=="Bob" )
then
    System.out.println($p);
end
```



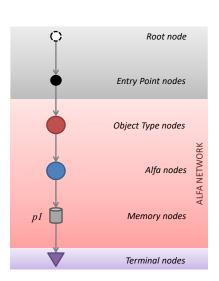
```
rule "Find Bobs"
when
  $p: Person( name=="Bob" )
then
  System.out.println($p);
end
    WM
```



```
rule "Find Bobs"
when
  $p: Person( name=="Bob" )
then
  System.out.println($p);
end
               p1: Person("Bob", null)
     WM
Person[Bob, <null>]
```

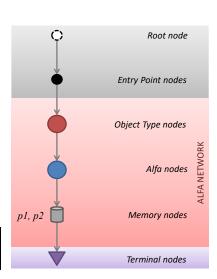


```
rule "Find Bobs"
when
  $p: Person( name=="Bob" )
then
  System.out.println($p);
end
                p1: Person("Bob", null)
                 a1: Address("Via Po 2", 40068,
                           "San Lazzaro")
     WM
Person[Bob, <null>]
```

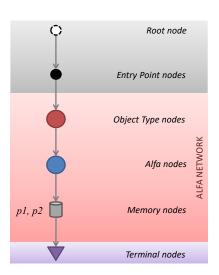


```
rule "Find Bobs"
when
  $p: Person( name=="Bob" )
then
  System.out.println($p);
end
                 p1: Person("Bob", null)
                 a1: Address("Via Po 2", 40068,
                           "San Lazzaro")
     WM
                 p2: Person("Bob", a1)
Person[Bob, <null>]
```

Person[Bob, Address[Via Po 2, 40068, San Lazzaro]]



```
rule "Find Bobs"
when
  $p: Person( name=="Bob" )
then
  System.out.println($p);
end
                  p1: Person("Bob", null)
                  a1: Address("Via Po 2", 40068,
                             "San Lazzaro")
      WM
                  p2: Person("Bob", a1)
                  p3: Person("Frank", a1)
Person[Bob, <null>]
Person[Bob, Address[Via Po 2, 40068, San Lazzaro]]
```



```
rule "Find Bobs and addresses"
when
$a: Address()
$p: Person( name=="Bob" )
then
System.out.println($p+"/"+$a+" ");
```

end

Root node

Entry Point nodes

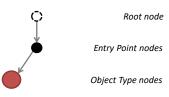
```
rule "Find Bobs and addresses"
when

$a: Address()

$p: Person( name=="Bob" )
then

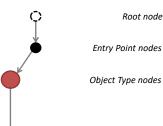
System.out.println($p+"/"+$a+" ");
```

end



```
rule "Find Bobs and addresses"
when

$a: Address()
$p: Person( name=="Bob" )
then
    System.out.println($p+"/"+$a+" ");
end
```



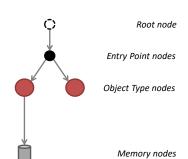
Memory nodes

```
rule "Find Bobs and addresses"
when

$a: Address()

$p: Person( name=="Bob" )
then

System.out.println($p+"/"+$a+" ");
end
```

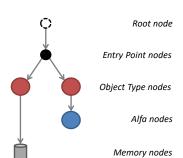


```
rule "Find Bobs and addresses"
when

$a: Address()

$p: Person(| hame=="Bob" )
then

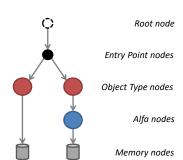
System.out.println($p+"/"+$a+" ");
end
```



```
rule "Find Bobs and addresses"
when

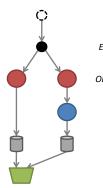
$a: Address()

$p: Person(| name=="Bob" )
then
    System.out.println($p+"/"+$a+" ");
end
```



```
rule "Find Bobs and addresses"
when

$a: Address()
$p: Person(| hame=="Bob" )
then
    System.out.println($p+"/"+$a+" ");
end
```



Root node

Entry Point nodes

Object Type nodes

Alfa nodes

Memory nodes

Beta nodes

```
rule "Find Bobs and addresses"
when

$\frac{\$a: Address()}{\$p: Person(|\name=="Bob")}
then
    System.out.println(\$p+"/"+\$a+" ");
end

Alfa nodes

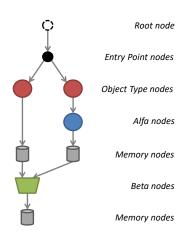
Memory nodes
```

NB: Beta Nodes make cartesian product of objects filtered by Alfa father!

Beta nodes

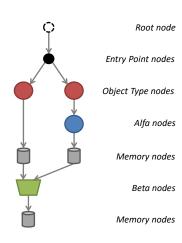
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when

$a: Address()
$p: Person(|name=="Bob" )
then
System.out.println($p+"/"+$a+" ");
end
```



```
rule "Find Bobs and addresses"
when

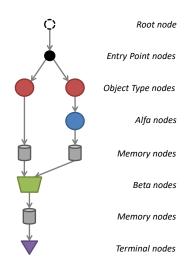
$a: Address()
$p: Person(|name=="Bob" )
then
System.out.println($p+"/"+$a+" ");
end
```



NB: tuple in a(Beta) Memory Node match with a composite pattern!

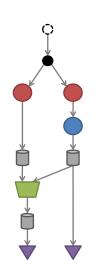
```
rule "Find Bobs and addresses"
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$a: Address()
$p: Person(|name=="Bob" )
then
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end
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```
rule "Find Bobs and addresses"
when

$a: Address()
$p: Person(|name=="Bob" )
then
System.out.println($p+"/"+$a+" ");
end
```



Root node

Entry Point nodes

Object Type nodes

Alfa nodes

Memory nodes

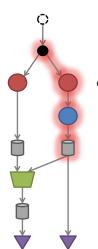
Beta nodes

Memory nodes

Terminal nodes

```
rule "Find Bobs and addresses"
when

$a: Address()
$p: Person(|hame=="Bob" )
then
System.out.println($p+"/"+$a+" ");
end
```



Root node

Entry Point nodes

Object Type nodes

Alfa nodes

Memory nodes

Beta nodes

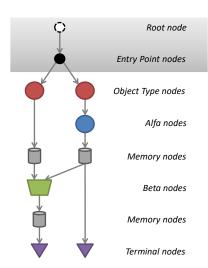
Memory nodes

Terminal nodes

NB: nodes of the previous rules are shared!

```
rule "Find Bobs and addresses"
when

$a: Address()
$p: Person( | name=="Bob" )
then
System.out.println($p+"/"+$a+" ");
end
```

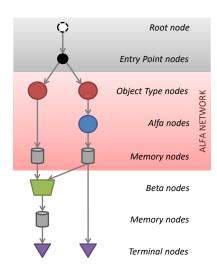


```
rule "Find Bobs and addresses"
when

$a: Address()
$p: Person(| hame=="Bob" )
then

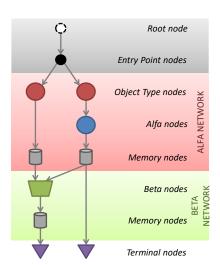
System.out.println($p+"/"+$a+" ");
```

end



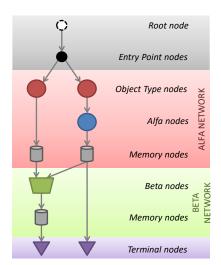
```
rule "Find Bobs and addresses"
when

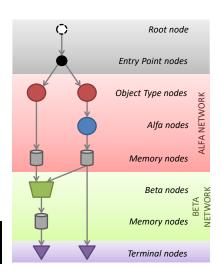
$a: Address()
$p: Person(|| hame=="Bob" )
then
    System.out.println($p+"/"+$a+" ");
end
```

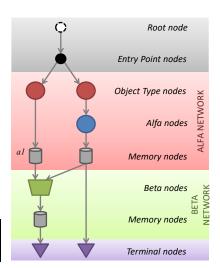


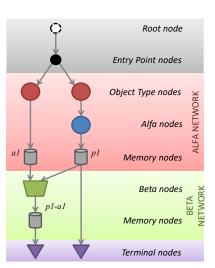
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$a: Address()
$p: Person(|name=="Bob" )
then
    System.out.println($p+"/"+$a+" ");
end
```



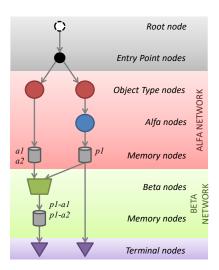






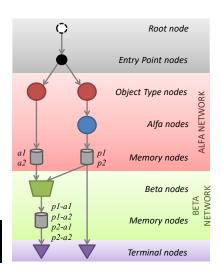
Person[p1, -]/Address[a1]

```
rule "Find Bobs and addresses"
when
  $a: Address()
  $p: Person( name=="Bob"
hen
  System.out.println(p+"/"+a+"");
end
                 a1: Address("Via Po 2", 40068,
                            "San Lazzaro")
                 p1: Person("Bob", null)
                 a2: Address("Via Roma 5",
     WM
                            40128, "Bologna")
```



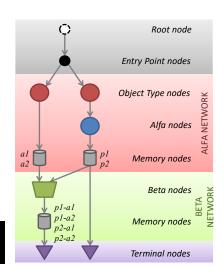
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when
  $a: Address()
  $p: Person( name=="Bob"
hen
  System.out.println(p+"/"+a+"");
end
                  a1: Address("Via Po 2", 40068,
                            "San Lazzaro")
                 p1: Person("Bob", null)
                  a2: Address("Via Roma 5",
     WM
                            40128, "Bologna")
                 p2: Person("Bob", a1)
```

Person[p2, -]/Address[a1] Person[p2, a1]/Address[a2]



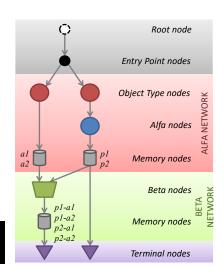
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  System.out.println(p+"/"+a+"");
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                  a1: Address("Via Po 2", 40068,
                            "San Lazzaro")
                  p1: Person("Bob", null)
     WM
                  a2: Address("Via Roma 5",
                             40128, "Bologna")
                  p2: Person("Bob", a1)
                  p3: Person("Giacomo", a1)
```

Person[p2, -]/Address[a1] Person[p2, a1]/Address[a2]



```
rule "Find Bobs and addresses"
when
  $a: Address()
  $p: Person( name=="Bob"
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                  a1: Address("Via Po 2", 40068,
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                  p1: Person("Bob", null)
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                  a2: Address("Via Roma 5",
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                  p2: Person("Bob", a1)
                  p3: Person("Giacomo", a1)
```

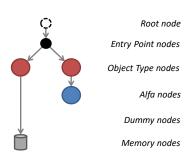
Person[p2, -]/Address[a1] Person[p2, a1]/Address[a2]

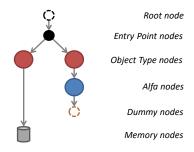


```
rule "Find Bob with its address"
when

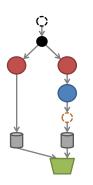
$a: Address()

$p: Person( | name=="Bob", | address == $a )
then
    System.out.println($p);
end
```





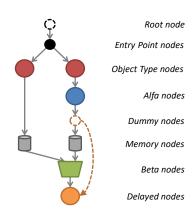
NB: this Alfa node contains a cross reference that cannot be resolved.



Root node
Entry Point nodes
Object Type nodes
Alfa nodes
Dummy nodes

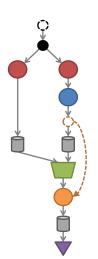
Memory nodes

Beta nodes



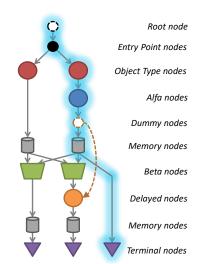
NB: the previous Alfa node is inserted here because it can resolve the cross reference.

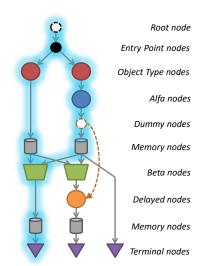
end

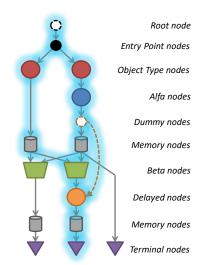


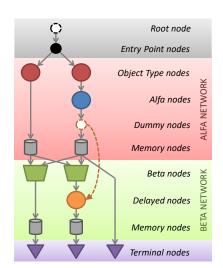
Root node **Entry Point nodes** Object Type nodes Alfa nodes Dummy nodes Memory nodes Beta nodes Delayed nodes Memory nodes

Terminal nodes

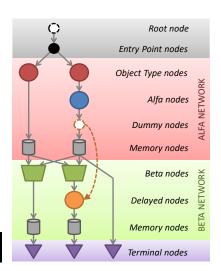








```
rule "Find Bob with its address"
when
  $a: Address()
  $p: Person( name=="Bob",
              address == $a )
then
  System.out.println($p);
end
     WM
```



```
rule "Find Bob with its address"

when

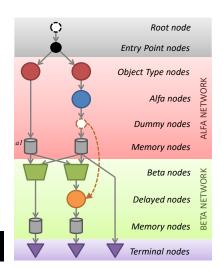
$a: Address()

$p: Person(||name==||Bob||, ||address|| == $a|)

then

System.out.println($p);
end

al: Address("Via Po 2", 40068, ||San Lazzaro")
```



```
rule "Find Bob with its address"

then

$a: Address()

$p: Person(| name=="Bob", address == $a|)

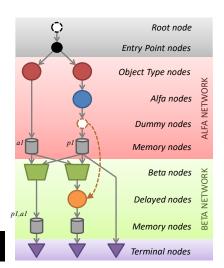
then

System.out.println($p);

end

al: Address("Via Po 2", 40068, "San Lazzaro")

pl: Person("Bob", null)
```



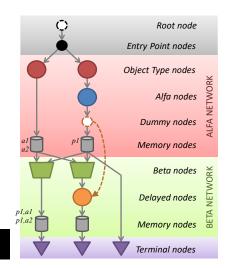
WM

"San Lazzaro")

40128, "Bologna")

p1: Person("Bob", null)

a2: Address("Via Roma 5",



end

WM

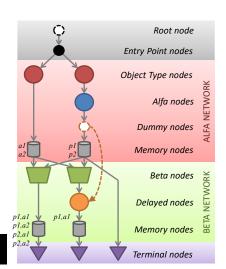
a1: Address("Via Po 2", 40068,

"San Lazzaro")
p1: Person("Bob", null)

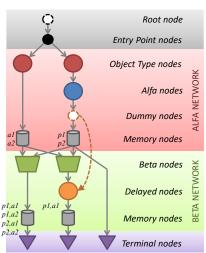
a2: Address("Via Roma 5",

40128, "Bologna") p2: Person("Bob", a1)

Person[Francesco, Address[Via Po 2, 40068, San Lazzaro]]

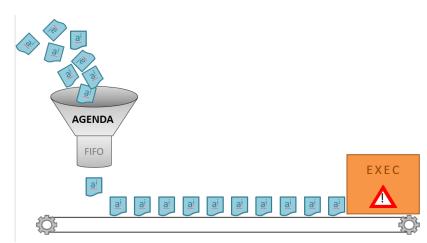


```
rule "Find Bob with its address"
when
  $a: Address()
  $p: Person( hame=="Bob",
                 address == $a )
then
  System.out.println($p);
end
                  a1: Address("Via Po 2", 40068,
                             "San Lazzaro")
                  p1: Person("Bob", null)
                  a2: Address("Via Roma 5".
      WM
                              40128, "Bologna")
                  p2: Person("Bob", a1)
                  p3: Person("Frank", a1)
Person[Francesco, Address[Via Po 2, 40068, San Lazzaro]]
```



CONFLICT RESOLUTION AND EXECUTION



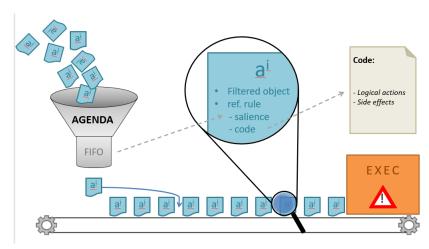


Rete00 Examples

CONFLICT RESOLUTION AND EXECUTION



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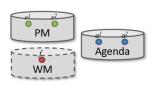


Conflict resolution and

Execution

CONFLICT RESOLUTION AND EXECUTION





```
rule "r1"
when
  F()
then
  assert(new G());
end
```

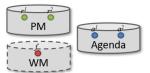




CONFLICT RESOLUTION AND EXECUTION



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First insert G, then retract F.

$a^2 < a^1$



First retract F, a1 cannot be applied, G never inserted.

$$r^1 < r^2$$

rule "r1" rule "r2" salience 10 salience 5

Estabilish a precedence fixed order between r1 and r2.

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